

UNITS 7 AND 8: SEMANTIC ANALYSIS and ERROR HANDLING

We want to incorporate a repetitive sentence into a high-level language. The sentence can be represented by the following regular expression:

repeat (identifier | number) >> sentence⁺ <<

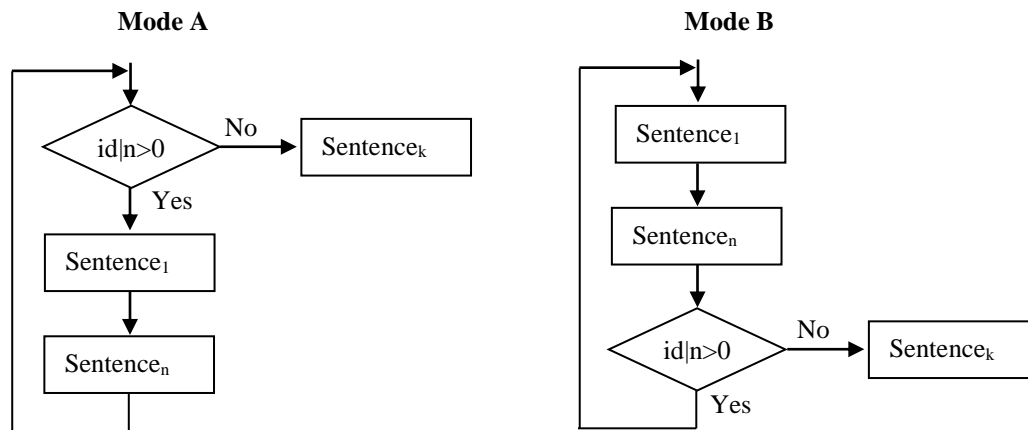
A program consists of at least one statement, where statements can be assignments, conditionals, and loops.

NOTE: The symbols "|" and "+" are part of the regular expressions, the others are part of the language.

It is required:

1. Define the grammar G that would generate valid programs of this programming language. Consider the assignment and conditional statements as terminal symbols of the grammar.
2. Describe the semantic routines of the grammar G that generate intermediate code in quartets with the following instructions, where pos are memory addresses, registers, or a number, and reg , $reg1$ and $reg2$ can be a record or a number. Write the semantic routines for the two interpretations that can be made about the execution flow of the loop:

```
repeat (id | n) >>  
sentence1  
...  
sentencen  
<<  
sentencek
```



Statement	Meaning
(move, pos_1, pos_2)	$pos_2 \leftarrow pos_1$
(push, $pos_1, ,$)	includes the contents of pos_1 into the stack
(pop, $, , pos_1$)	$pos_1 \leftarrow$ top of the stack
(label, $, , label$)	defines a label
(goto, $, , label$)	go to a label
(return, $, , reg$)	go to the address specified by reg
(if, $reg, , label$)	go to label reg es -1
(<, $reg, , label$)	go to label if the contents of reg is lower or equal to 0
(+, reg_1, reg_2, reg)	$reg \leftarrow reg_1 + reg_2$
(-, reg_1, reg_2, reg)	$reg \leftarrow reg_1 - reg_2$
(*, reg_1, reg_2, reg)	$reg \leftarrow reg_1 * reg_2$
(/, reg_1, reg_2, reg)	$reg \leftarrow reg_1 / reg_2$

SOLUTION:

A grammar to generate the language defined:

$G = \{ \text{assignment, condition, id, n, repeat, (,), <<, >> \}, \{ \mathbf{S}, \mathbf{S}', \mathbf{B}, \mathbf{E}, \mathbf{R} \}, \{ \mathbf{S} \}$

- (1) $\mathbf{S} ::= \mathbf{E} \mathbf{S}'$
- (2) $\mathbf{S}' ::= \mathbf{S}$
- (3) $\mathbf{S}' ::= \lambda$
- (4) $\mathbf{E} ::= \text{assignment}$
- (5) $\mathbf{E} ::= \text{condition}$
- (6) $\mathbf{E} ::= \mathbf{B}$
- (7) $\mathbf{B} ::= \text{repeat} (\mathbf{R}$
- (8) $\mathbf{R} ::= \text{id}) >> \mathbf{S} <<$
- (9) $\mathbf{R} ::= \text{n}) >> \mathbf{S} <<$

O::=Id

O.Value:= Id.Value;
O.Code:= "";

O::=Num

O.Value:= Id.Value;
O.Code:= "";

Z::=E

Z.Value:=E.Value;
Z.Code:=E.Code;

Z::=C

Z.Value:=C.Value;
Z.Code:=C.Code;

S::=λ

S.Code:= ""

S₀::=CS₁

S₀.Code:=C.Code
S₁.Code

S₀::=CS₁

S₀.Code:=E.Code
S₁.Code

Use the stack to know which variables to assign the value of the expression, stack = -1 empty.

T::=λ

T.Code:=(push,-1,,)

T::=V

T.Code:=V.Code

V::= Id T

V.Value:=newtemp;
V.Code:=(push,,,Id)

O::=+

O'.Code=""
O'.Operation="+"

The Operation attribute is included to later know which operation to perform.

O'::=-

O'.Code=""
O'.Operation="-"

O'::=*

O'.Code=""
O'.Operation="*"

O'::=/
O'.Code=""
O'.Operation="/"

U::=λ

U.Code:= ""

U::=O' E'

U.Value:=E'.Value;
U.Operation=O'.Operation
U.Code:=E'.Code

E'::=O U

E'.Value:=newtemp;
E'.Code:= O.Code
if U.Code=""
then (move, O.Value, , E'.Value)
else U.Code
Select case U.Operation
case "+"
(move, O.Value,,A)
(move, U.Value,,B)
(+, O.Value,,A)
case "-"
(move, O.Value,,A)
(move, U.Value,,B)
(-, O.Value,,A)
case "*"
(move, O.Value,,A)
(move, U.Value,,B)
(*, O.Value,,A)
case "/"
(move, O.Value,,A)
(move, U.Value,,B)
(/, O.Value,,A)
end select
(move, A,, E'.Value)

E'::=E'->V

E.Start:=newlabel;
E.Stack_Empty:=newlabel;
E.Code:=(label,,,Start)
(pop,,,A)
(if,A,,Stack_Empty)
(move,A,E'.Value.)
(goto,,,Start)
(label,,,Stack_Empty)

C::=,(E')=>ZW

C.False:=newlabel;
C.Exit:=newlabel;
C.Code:=E'.Code
(<,E'.Value,,C.False)
Z.Code
(goto,,,C.Exit)
(label,,,C.False)
W.Code
(label,,,C.Exit)

W::=?

W.Code:= ""

W::=/=>Z?

W.Code:= Z.Code