



# ***Introduction to digital systems and microprocessors***

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# Outline

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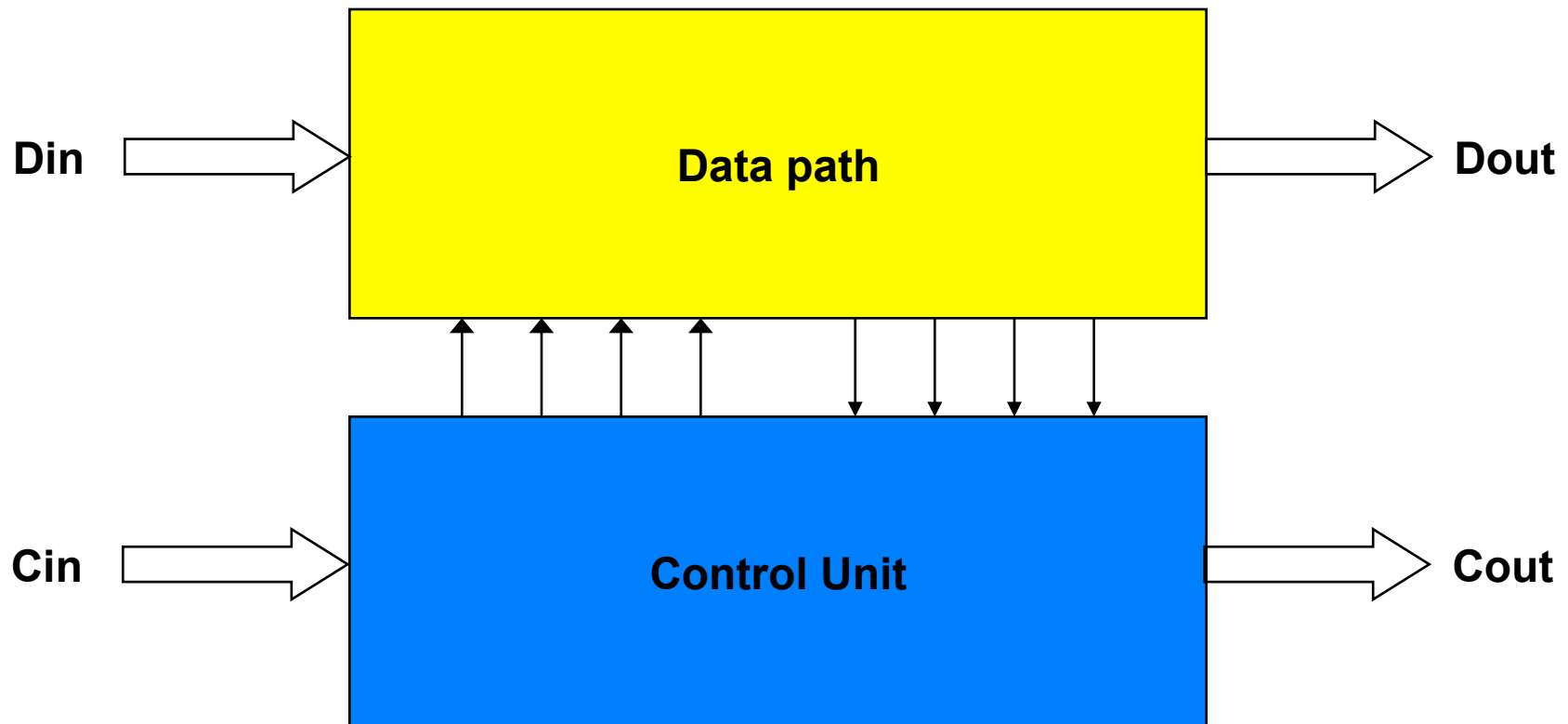
- Structure of a digital system
  - Data path
  - Control unit
- Structure of a basic computer
- Operation of a basic computer. Instructions

# Digital systems

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- A digital system processes digital information according to a particular algorithm
- Algorithm: a finite sequence of operations to solve a problem
- Goal of this lesson: introduce digital systems

# Structure of a digital system



# The Datapath

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- Any algorithm can be decomposed in a set of basic operations:
  - Arithmetic
  - Logic
  - Shift & Rotation
- The datapath is the set of functional units that process the data

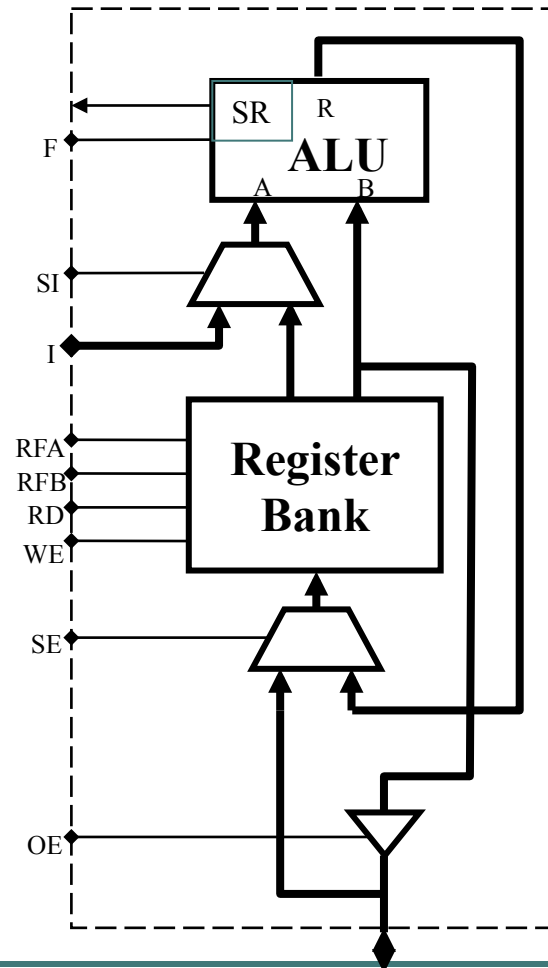
# The Datapath

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- Typical datapath components:
  - ALUs: perform operations
  - Registers and memories: store temporal data
  - Buses: connect datapath elements
  - Multiplexers: select data to be processed every time

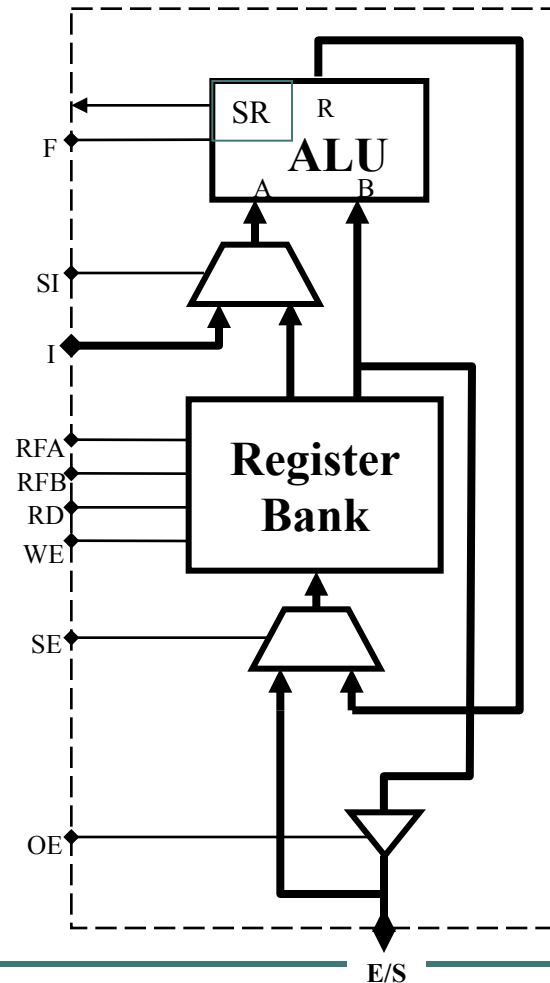
# Datapath example

- ALU performs operations
  - F: ALU function
- SR (Status Register). Flag bits resulting from the ALU operation, such as :
  - C: Carry
  - O: Overflow
  - Z: Zero
  - S: Sign
- Immediate operands provided by I and selected by SI



# Datapath example

- Register bank (triple port memory) stores temporal data
- Register selection
  - RFA: Source register A
  - RFB: Source register B
  - RD: Destination register
- I/O for external connection
  - SE selects RD input (external or internal)
  - OE (Output Enable)





# The Control Unit

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- Determines the correct application and sequencing of operations on data
- Typical Control Unit components:
  - Finite State Machines
  - Counters
  - Registers and memories to store control data, etc.

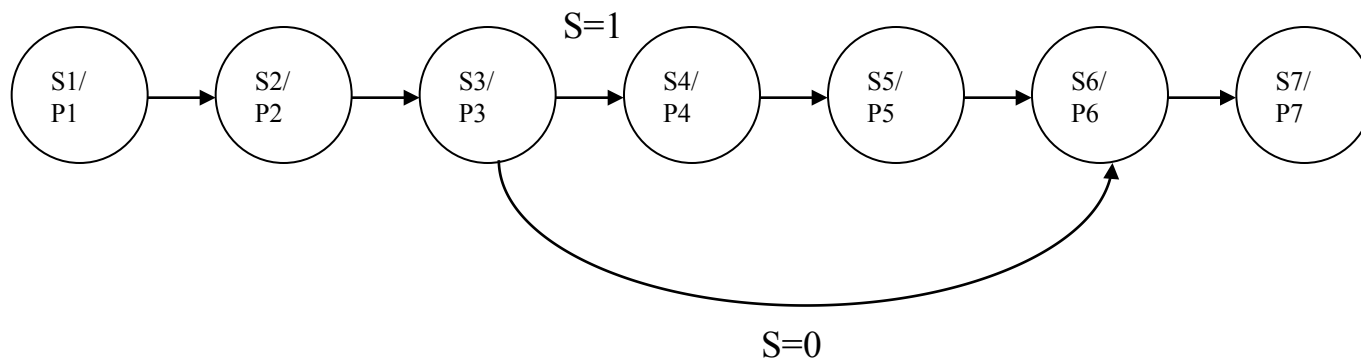
# Control Unit example

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- Example: perform the following computation
$$y = \text{abs}(x1 - x2)/2$$
- Steps:
  - 1. Load  $x1$  in register R1
  - 2. Load  $x2$  in register R2
  - 3. Subtract  $x2$  from  $x1$  and place the result in register R3 ( $R3 = R1 - R2$ ).
    - If the result is positive, go to step 6. Otherwise, continue with step 4.
  - 4. Complement R3 ( $R3 = \text{NOT } R3$ )
  - 5. Increment R3 ( $R3 = R3 + 1$ )
  - 6. Shift right R3 ( $R3 = R3/2$ )
  - 7. Send R3 to the output

# Control Unit example

- FSM of the Control Unit



# Control Unit example

- FSM outputs

Paso	A L U Function (F)	RFA	SI	I	RFB	RD	WE	SE	OE
1	X	X	X	X	X	R1	0	1	0
2	X	X	X	X	X	R2	0	1	0
3	RESTA	R1	0	X	R2	R3	0	0	0
4	NOT	R3	0	X	X	R3	0	0	0
5	SUMA	X	1	1	R3	R3	0	0	0
6	SHR	R3	0	X	X	R3	0	0	0
7	X	X	X	X	R3	X	1	X	1

# Digital systems design

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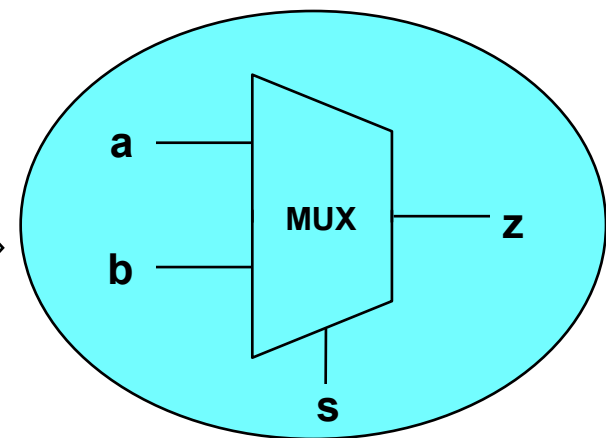
- Designing with very basic components, such as logic gates, is not practical for digital systems design
- RTL (Register Transfer Level) design:
  - Use more abstract components, such as ALUs, registers or multiplexers
  - Digital systems are described by means of operations performed with data stored in registers

# Hardware Description Languages

- Hardware Description Languages (HDLs) allow a designer to work at a higher abstraction level
- Example (VHDL)

```
if s = '0' then  
    z <= a;  
else  
    z <= b;  
end if;
```

Synthesizer  
(Compiler)



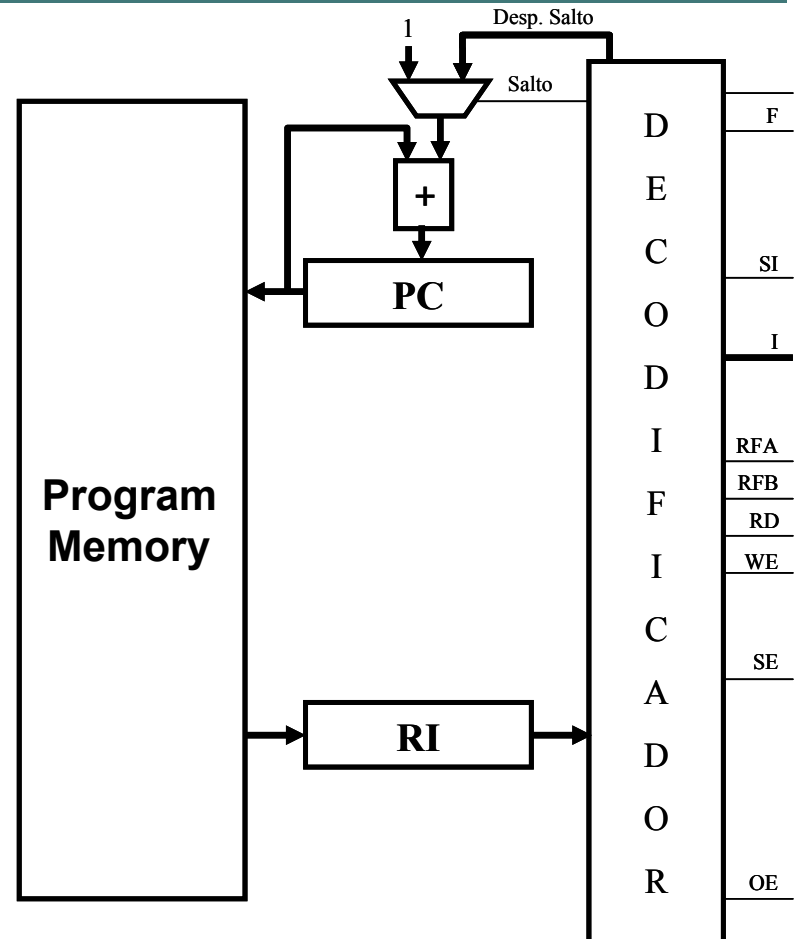
# Structure of a basic computer

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- Computer: general purpose machine for information processing
  - Oriented to implement a wide variety of algorithms
- Components:
  - Processing Unit or Central Processing Unit (CPU) performs the required operations
    - General purpose datapath
    - Programmable control unit, able to perform many different algorithms
  - Memory to store information
  - Input/Output units for external communication

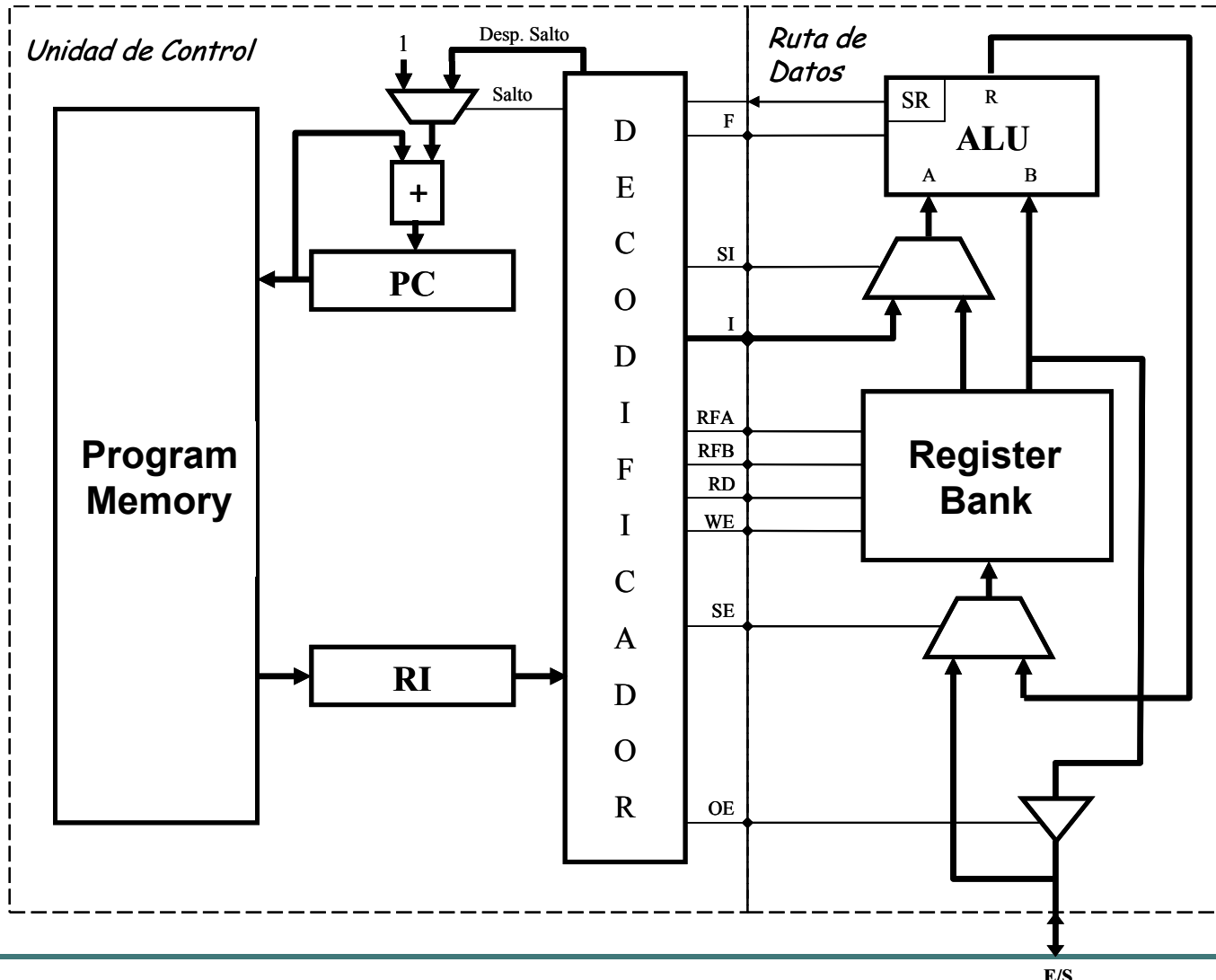
# Programmable Control Unit

- Control words are stored in a memory
  - Control words are coded to save space
  - Coded control words are called *instructions*
- Every clock cycle, an instruction is read (RI) and decoded
- The instruction required for each step is determined by the Program Counter (PC).
  - Sequence changes are implemented by means of instructions, that can dynamically depend on the SR value

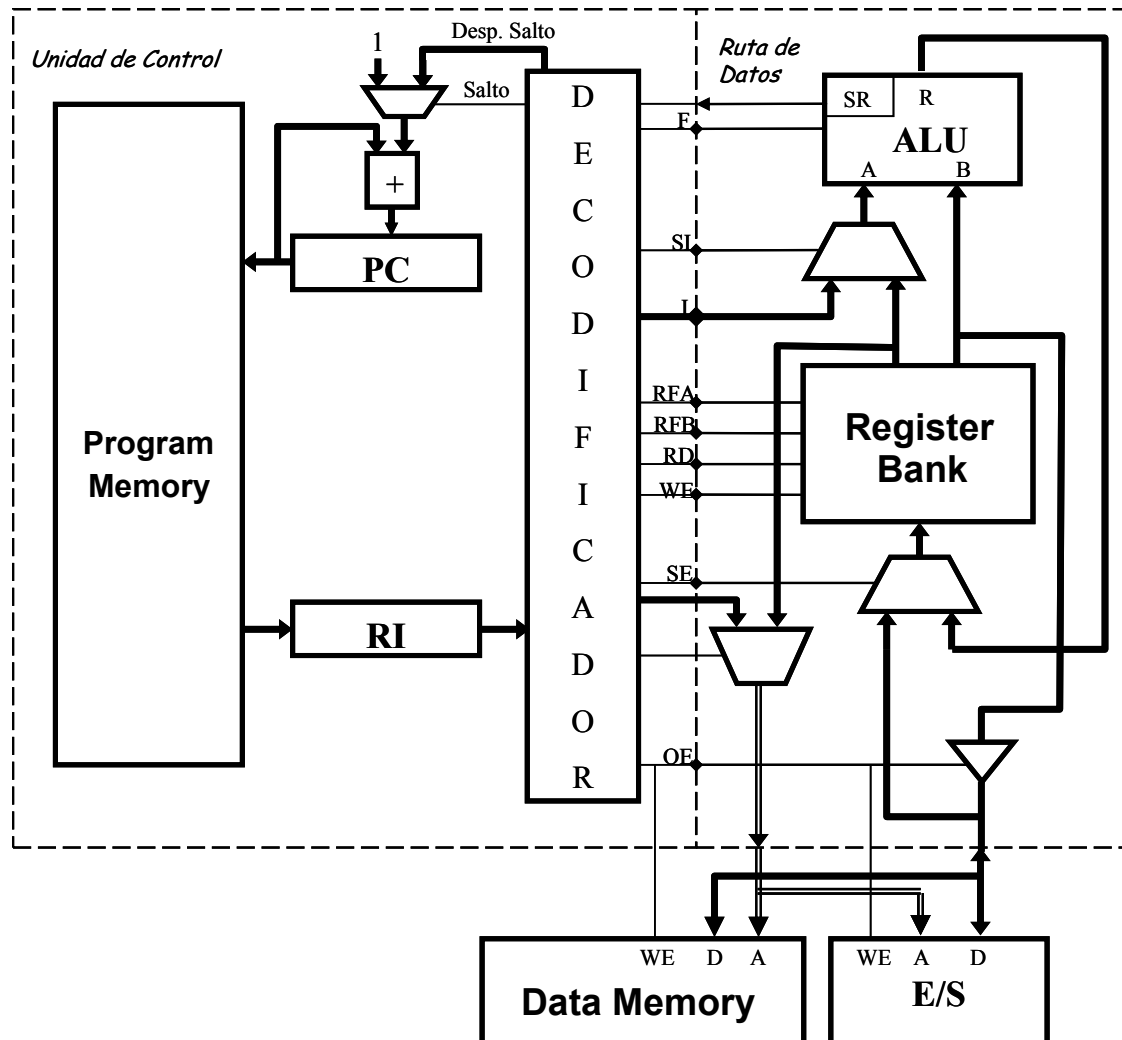




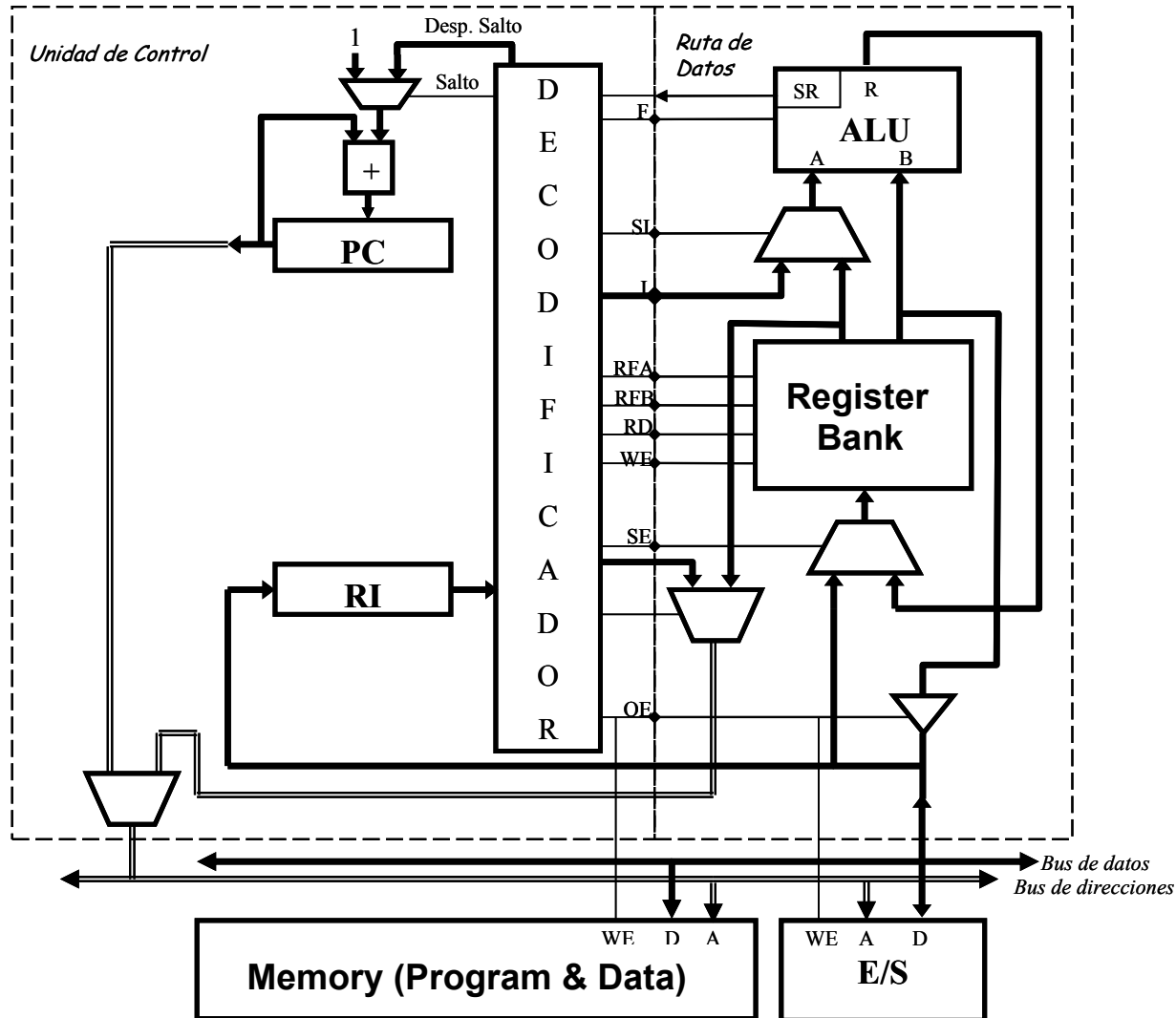
# Microprocessor



# Basic computer (Harvard architecture)



# Basic computer (von Neuman architecture)



# Operation of the basic computer.

## Instructions

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- How are instructions executed?
  - Instruction cycle
- What are the main types of instructions?
- How are instructions coded?
  - Instruction format
  - Addressing modes

# Instruction cycle

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- Every instruction is typically processed in two steps:
  - Instruction fetch: load instruction in IR
  - Instruction execution: decode instruction and configure the datapath to execute the operation specified by the instruction
- Each of these steps can be executed in a single clock cycle or in many clock cycles, depending on the microprocessor's complexity

# Types of instructions

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- Data transfer instructions
  - Transfer data between registers, register and memory, or register and I/O
- Arithmetic and logical instructions
  - Set ALU operations: add, subtract, shift, etc.
- Control flow instructions
  - Modify the instruction execution sequence
  - Modify PC contents
  - Examples: branches, subroutine calls, etc.

# Instruction format

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- Organization of instruction information. Instructions are typically field coded:
  - Operation code (Opcode): specifies the operation to be performed
  - Operands: specifies the data to be operated
- The number of operands and the size of each field may vary
- Instruction size must be an integer multiple of memory width

# Addressing modes

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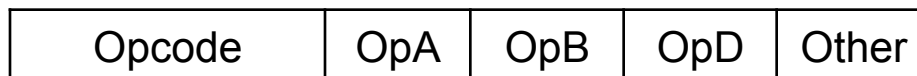
- Operands can be specified in different ways, called addressing modes
- Some usual addressing modes:
  - *Immediate*: the operand is given with the instruction
  - *Register direct*: the instruction gives a register that contains the operand
  - *Memory direct*: the instruction gives a memory address that contains the operand
  - *Indirect*: the instruction gives a register that contains the memory address of the operand



# Instruction format example

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- Consider the previous microprocessor architecture, with the following parameters:
  - Opcode field: 5 bits
  - Register bank: 8 registers (direct addressing with 3 bits)
- Instruction format



# Instruction examples

O p e r a t i o n	Mnemonic	Operation Code
Load data from memory	L D	0 0 0 0 0
Store data in memory	S T	0 0 0 0 1
A d d	A D D	0 1 0 0 0
S u b t r a c t	S U B	0 1 0 0 1
N O T	N O T	0 1 1 0 0
A N D	A N D	0 1 1 0 1
O R	O R	0 1 1 1 0
S h i f t	SHL, SHR	0 1 1 1 1
J u m p	J M P	1 0 0 0 0
Subroutine call	C A L L	1 0 1 0 0
Subroutine return	R E T	1 0 1 0 1
.	.	.

# Instruction examples

- Single word instruction

Opcode	RFA	RFB	RD	Other
01000	001	010	011	00

**R3 = R1 + R2**

- Double word instruction

Opcode	RFA	RFB	RD	Other
00000	000	000	011	00
Direccion				
1010 1011 1100 1101				

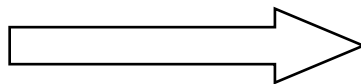
**Load R3 with data in memory position ABCDh**

# Assembler language

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- Binary instruction codes are cumbersome
- *Assembly Language*: instructions are specified as mnemonics and operands are specified using symbolic names
- Assembly Program: translate instructions in a program to their equivalent binary codes
- Example:

**ADD R1, R2, R3**



0100000101001100

# References

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- “Principios de Diseño Digital”. D. Gajski. Ed. Prentice Hall