## uc3m Universidad Carlos III de Madrid

DEPARTMENT OF COMPUTER SCIENCE CARLOS III UNIVERSITY OF MADRID

**Computer Science** Language Processors

## Rules

- The duration of the test is **60 minutes**
- Questions will not be answered during the test
- One cannot re-enter the classroom after leaving it
- The answers must be written using a pen (not a pencil)

## 1.- Given the grammar:

D ::= bA | cX | d X ::= MbA | BF M ::= cM | 1 A ::= MBj | F B ::= c | 1 F ::= fA | 1

a) Calculate the FIRST and FOLLOW sets.

 $FIRST(D) = \{b, c, d\}$   $FIRST(X) = \{b, c, f, \lambda\}$   $FIRST(M) = \{c, \lambda\}$   $FIRST(A) = \{c, f, j, \lambda\}$   $FIRST(B) = \{c, \lambda\}$  $FIRST(F) = \{f, \lambda\}$ 

FOLLOW(D) = {\$} FOLLOW(X) = {\$} FOLLOW (M) = {b, c, j} FOLLOW (A) = {\$} FOLLOW (B) = {j, f, \$} FOLLOW (F) = {\$}

#### b) Using the algorithm, determine if it is an LL(1) grammar.

For a grammar to be a LL(1) grammar, it must fulfill that there are not two or more productions in any cell of the analysis table. This condition will occur when:

 $\forall$  production A ::=  $\alpha_i$ |...| $\alpha_n$ :

- FIRST( $\alpha$  i)  $\cap$  FIRST( $\alpha$  i) = 0  $\forall$  i  $\neq$  j
- If  $\alpha i ::= \lambda$  then FIRST( $\alpha i$ )  $\cap$  FOLLOW(A) = 0  $\forall i \neq j$



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For the given grammar:

FIRST (bA)  $\cap$  FIRST(cX)  $\cap$  FIRST(d) = 0 FIRST (MbA)  $\cap$  FIRST(BF) = {c, b}  $\cap$  {c, f,  $\cap$ }  $\cap$  0 FIRST (cM)  $\cap$  FIRST( $\lambda$ ) = 0 FIRST (cM)  $\cap$  FOLLOW(M) = {c}  $\cap$  {c, j}  $\cap$  0 FIRST (MBj)  $\cap$  FIRST(F) = {c,  $\cap$ }  $\cap$  {f,  $\cap$ }  $\cap$  0 FIRST (c)  $\cap$  FIRST( $\lambda$ ) = 0 FIRST (c)  $\cap$  FOLLOW(B) = {c}  $\cap$  {j, f, \$} = 0 FIRST (fA)  $\cap$  FOLLOW( $\lambda$ ) = 0 FIRST (fA)  $\cap$  FOLLOW(F) = {f}  $\cap$  {\$} = 0

Then, the grammar is not an LL(1) grammar.

## c) Construct the analysis table for the LL(1) table-driven top-down predictive parsing.

The parsing table for the grammar is:

	b	с	d	j	f	\$
D	D → bA	$D \rightarrow cX$	$D \rightarrow d$			
Χ	X → MbA	X → MbA			$X \rightarrow BF$	$X \rightarrow BF$
		$X \rightarrow BF$				
Μ	$M \rightarrow \lambda$	$M \rightarrow cM$		$M \rightarrow \lambda$		
		$M \rightarrow \lambda$				
Α		A → MBj		A → MBj	$A \rightarrow F$	$A \rightarrow F$
B		$B \rightarrow c$		$B \rightarrow \lambda$	$B \rightarrow \lambda$	$B \rightarrow \lambda$
F					$F \rightarrow fA$	$F \rightarrow \lambda$



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