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# Lesson 7 Algorithms with arrays

### Programming

Grade in Computer Engineering





- 1. Search
  - Linear

**Binary** 

- 2. Sort
  - Bubble
  - Selection
  - Insertion



#### Outline

1. Search Linear

## Binary

- 2. Sort
  - Bubble
  - Selection
  - Insertion



Search algorithms aim at finding a value in a collection (usually, the first occurrence) Input: Array of values *list*, value to find *e* Output: Position of *e* in *a*; -1 if not found

```
Find a value in an array of integers
public static int find(int [] list, int e)
    list = {5, 6, 3, 1, 8, 9, 0, 2, 4, 1, 7}
    e = 1
Output:
```

3



## Looks for the value *e* sequentially in *list*:

```
location = -1;
i = 0;
found = false;
while ( (!found) && (i < list.length))
if (list[i] == e)
    location = i;
    found = true;
    i++;
return location;
```



## At most, *List.Length* tests are needed Without further assumptions on *List*, it is the most efficient (non-parallel) search algorithm

```
/** Linear search
 * @param list Sequence of elements
 * @param e Element to find
 * @return Position of e in list; -1 if not found */
public static int linearSearch(int [] list, int e) {
    int location = -1;
    boolean found = false;
    for(int i=0; i<list.length && !found; i++)
        if(list[i] == e) {
            location = i;
            found = true;
        }
    return location;
}
search.basic.Algorithms</pre>
```



If *List* is ordered, we can narrow the search to one half of the array:

```
location = -1;
left = 0;
right = list.length - 1;
middle = list.length / 2;
found = false;
while ( (left <= right) and (!found) )</pre>
    if ( list[middle] == e )
         found = true;
         location = middle;
    else if (e < list[middle] )</pre>
         right = middle - 1;
    else
         left = middle + 1;
    middle = (left+right) / 2;
return location;
```



#### At most, Log(List.Length) comparison are needed

```
/** Binary search
 * @param list Ordered sequence of elements
 * @param e Element to find
 * @return Position of e in list; -1 if not found */
public static int binarySearch(int [] list, int e) {
    int location = -1;
    int left = 0;
    int right = list.length - 1;
    int middle = list.length / 2;
    boolean found = false;
    while(left <= right && !found) {</pre>
        if(e == list[middle]) {
            found = true;
            location = middle;
        } else if(e < list[middle]) {</pre>
            right = middle-1;
        } else {
            left = middle+1;
        3
        middle = (right+left) / 2;
    }
                                                     search.basic.Algorithms
    return location;
}
```



#### Outline

## 1. Search Linear Binary

2. Sort Bubble Selection Insertion



Sort algorithms aim at rearranging the values of a collection to position them in order (usually, in increasing order)

- Input: Array of values list
- Output: Array of values *list\** ordered

```
Sort array
public static void sort(int [] list)
    list = {5, 6, 3, 1, 8, 9, 0, 2, 4, 1, 7}
Output:
```

```
list = \{0, 1, 1, 2, 3, 4, 5, 6, 7, 8, 9\}
```



### 2. Sorting algorithms Algorithms (direct)

## **Direct sorting algorithms**

- Most popular
  - Direct Swapping (Bubble sort)
  - Direct Insertion (Insertion sort)
  - Direct Selection (Selection sort)

#### Features

- Simple algorithms
- Not very efficient: complexity is  $\mathcal{O}(n^2)$
- Can be used with small arrays



## **Advanced sorting algorithms**

- Most popular
  - Shell
  - Quicksort
  - Heapsort
- Features
  - Sophisticated algorithms
  - Efficient: complexity is  $\mathcal{O}(n^* \log n)$
  - Are used with large arrays

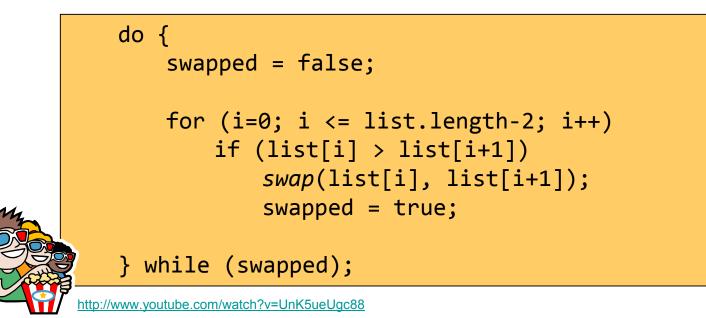
#### **Just for fun!** Bogosort: Random reordering of the array

H. Gruber, M. Holzer and O. Ruepp: <u>Sorting the Slow Way: An Analysis of Perversely Awful Randomized Sorting Algorithms</u>, 4th International Conference on Fun with Algorithms, Castiglioncello, Italy, 2007, Lecture Notes in Computer Science 4475, pp. 183-197.



## Idea:

- Compare an element *List[i]* with the adjacent value *List[i+1]*
- If List[i] > List[i+1], the values are swapped
- Repeat the procedure for the complete array while swaps are performed





```
/** Bubble sort
 * @param list Array to sort */
public static void bubbleSort(int [] list) {
    boolean swapped;
    do {
        swapped = false;
        for(int i=0; i <= list.length-2; i++) {</pre>
            if(list[i] > list[i+1]) {
                 int temp = list[i];
                 list[i] = list[i+1];
                 list[i+1] = temp;
                 swapped = true;
            }
        }
    } while(swapped);
}
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```



Worst case: The array is reversed  $\mathcal{O}(n^2)$ The outer while is executed *n* times, since swapping is always performed

**Best case:** Array is ordered  $\mathcal{O}(\mathbf{n})$ swapped is not changed from false to true

#### Average: $\mathcal{O}(n^2)$

Swapping

Comparison

```
/** Bubble sort [.]
public static void bubbleSort(int [] list) {
    boolean swapped;
    do {
        swapped = false;
        for(int i=0; i <= list.length-2; i++) {</pre>
            if (list[i] > list[i+1]) {
                 int temp = list[i];
                list[i] = list[i+1];
                 list[i+1] = temp;
                 swapped = true;
            }
    } while(swapped);
```

}



## Idea:

For each value of the list (at position i),

Finds the smallest value (at position *minPos*) of the elements *i*+1,..., *List.Length-1* 

If List[i] > List[minPos], the values are swapped





```
/** Selection sort [.]
 public static void selectionSort(int [] list) {
     for(int i=0; i <= list.length-2; i++) {</pre>
         int minPos = i;
          for(int j=i+1; j < list.length; j++)</pre>
              if(list[j] < list[minPos])</pre>
                  minPos = j;
         int temp = list[i];
         list[i] = list[minPos];
         list[minPos] = temp;
     }
 }
```

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- The number of comparison operations does not depend on the initial order of the values. It will be equal to the number of evaluations of the condition of the *if* O(n<sup>2</sup>)
- The number of swap-related operations depends on the initial order of the values

Sw	appi	ng		
Со	mpa	riso	n	

```
/** Selection sort []
public static void selectionSort(int [] list) {
  for(int i=0; i <= list.length-2; i++) {
    int minPos = i;
    for(int j=i+1; j < list.length; j++)
        if(list[j] < list.length; j++)
        if(list[j] < list[minPos])
        minPos = j;
    int temp = list[i];
    list[i] = list[minPos];
    list[minPos] = temp;
  }
}</pre>
```



## Idea:

Assumes that the elements  $\theta$ , ..., *i-1* of the list are ordered

Finds the position k in 0, ..., i-1 where the element at position i should be placed

(Simultaneously) Shift to the right the values at k, ..., i-1 and inserts List[i] at position k

```
for (i=1; i < list.length; i++)
    e = list[i];
    j = i-1;
    while( (j >= 0) && (list[j] > e) )
        list[j+1] = list[j];
        j = j-1;
    list[j+1] = e;
```



```
/** Insertion sort ...
public static void insertionSort(int [] list) {
    for(int i=1; i < list.length; i++) {</pre>
        int e = list[i];
        int j = i-1;
        while(j \ge 0 \& list[j] > e) {
             list[j+1] = list[j];
             j--;
         }
        list[j+1] = e;
    }
}
```

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Worst case: The array is reversed 𝒪(n<sup>2</sup>)

The inner while is executed until *j* < 0 (max. number of iterations)

**Best case:** Array is ordered  $\mathcal{O}(n)$ The inner while is never executed

The inner while is never exec

Average:  $\mathcal{O}(n^2)$ 



Comparison

#### /\*\* Insertion sort ..

public static void insertionSort(int [] list) {

```
for(int i=1; i < list.length; i++) {
    int e = list[i];
    int j = i-1;
    while(j>=0 && list[j] > e) {
        list[j+1] = list[j];
        j--;
    }
    list[j+1] = e;
}
```

}



Algorithms can be compared according to the number of comparisons performed in the best case, worst case, and average case

Being *n* the length of the array:

Algorithm	Best ≈	Worst ≈	Average ≈
Bubble	n	n²	n²
Selection	n <sup>2</sup>	n <sup>2</sup>	n <sup>2</sup>
Insertion	n	n²	n <sup>2</sup>
Quicksort	n · log(n)	n²	n · log(n)



**Bubble sort** is the simplest, but also has a the higher worst-case execution time. Nevertheless, it behaves quite well with ordered arrays

*Selection sort* is easy to implement and more efficient that *Bubblesort*, but it behaves very bad even if the array is ordered (it cannot be known if the array is already sort at any iteration)

*Insertion sort* is simple to implement and behaves quite well for almost ordered arrays. It is also more efficient in practice



Develop a program to test the execution time of the three basic sorting methods for different array sizes = {1000, 2000, ..., 20000}

- The program must run 5 times each algorithm for an array size with different initial values.
- The program must generate three text files (*bubble.txt, selection.txt, insertion.txt*) with this structure:
  - <array size> <average> <best time> <worst time>
  - <array size> <average> <best time> <worst time>

•••

Represent the results (array size vs. average time) in a table and graphically (use Microsoft Excel).



### 2. Sorting algorithms Results

bubble.txt 🖾	
1000       4833840.0       3403156       8024134         2000       1.4096526E7       13973424       14377736         3000       3.1497941E7       31175959       32105338         4000       5.5688851E7       54822051       56         5000       8.6352455E7       85646500       86         6000       1.24527935E8       123430771       1000       7630089.0       442934       8024134         7000       1.69686803E8       168417844       3000       4.5609526E7       3918934       3210         4000       8.0775307E7       7016045       56866	7736 05338

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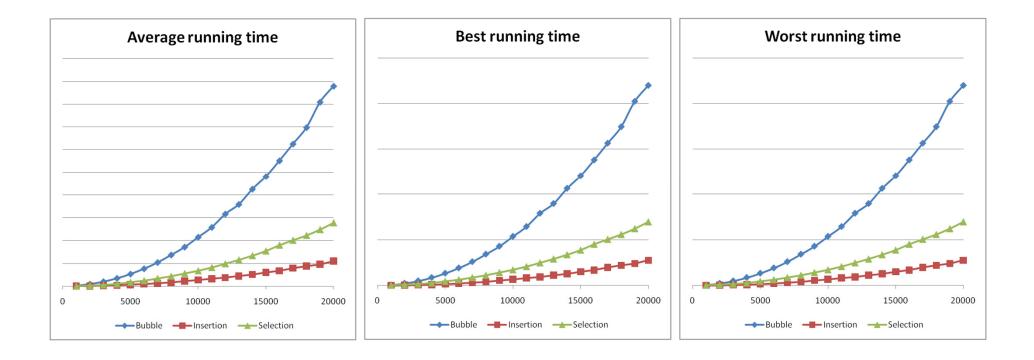


#### 2. Sorting algorithms Graphical respresentation

	Α	В	С	D	E	F	G	Н	1	J	k	(	L	-
1	Size	Average	Best	Worst								_		
2	1000	4833840	3403156	8024134										
3	2000	1,41E+07	13973424	14377736				Avera	ge					
4	3000	3,15E+07	31175959	32105338		1,6E+09 -								
5	4000	5,57E+07	54822051	56861696										
6	5000	8,64E+07	85646500	86805655		1,4E+09 -					$\neq$			
7	6000	1,25E+08	123430771	125393660						1	6			
8	7000	1,70E+08	168417844	170736644		1,2E+09 -				*				
9	8000	2,29E+08	224974962	238769942		1E+09 -								
10	9000	2,84E+08	281557013	287583059		12+09								
11	10000	3,49E+08	343487511	355747379		80000000								=
12	11000	4,24E+08	421367031	430040899										
13	12000	4,98E+08	496283885	501597131		60000000 -								
14	13000	5,85E+08	579794918	588853542					*					
15	14000	6,78E+08	675493952	680047464		40000000 -			*					
16	15000	7,77E+08	773596854	781549610		200000000		**						
17	16000	8,85E+08	882779445	889809180		200000000 -								
18	17000	1,01E+09	1001417149	1022621241		0 -		*						
19	18000	1,16E+09	1143381212	1199534019		-	50 50	000 1	0000 1	15000	20000			
20	19000	1,29E+09	1285006475	1310799766				_						
21	20000	1,46E+09	1434090715	1473585120			<ul> <li>Aver</li> </ul>	rage — Po	olinómica (Aver	rage)				
22												-		
23														-
14 -		Bubble /	Insertion	Selection / A	L <u>(</u> 2)			14						▶ [

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Results for 5 executions with random values in [0, 10)

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- 1. Search
  - Linear

**Binary** 

- 2. Sort
  - Bubble
  - Selection
  - Insertion



## Search

- Linear search
- **Binary search**

Use? Binary search if values are sorted; otherwise, linear search

### Sort

- **Bubble sort**
- Selection sort
- Insertion sort
- Use? None of them, go for Quicksort



## **Recommended lectures**

H. M. Deitel, P. J. Deitel. *Java: How to Program. Prentice Hall,* 2011 (9th Edition), Chapter 19 [link]



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