



AUTOMATA THEORY AND FORMAL LANGUAGES

UNIT 8: COMPUTATIONAL COMPLEXITY





OUTLINE

- Theory of Computation and Computational Complexity Theory
- Computational Models. Turing Machine
- Non-Deterministic Turing Machines
- Classification of problems
- Summary





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■ Theory of Computation:

- Branch of the Theoretical Computer Science that deals with whether and how efficiently problems can be solved on a model of computation, using an algorithm.
- It studies which problems are decidable using different formal Computation Models (Turing Machines, Lambda calculus, recursive functions, etc).





■ Computational Complexity Theory:

- Branch of the Theory of Computation focused on classifying computational problems according to their inherent difficulty.
- It studies the complexity level of the algorithm required to solve a decidable problem.





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■ Utilities of a Turing Machine

- Formalizes the concept of algorithm.
- It is a formal model of a computer, i.e.
 - “The recognition of the strings that make up a language is a formal way to express any problem.”
- Theoretical *human* computer.
- Provides means to prove:
 1. Whether **there is an algorithm** or procedure to solve a specific problem
 2. The **form/complexity** of the algorithm
 3. Time spent by this algorithm or procedure to solve the problem.





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Non-Deterministic Turing Machine

- They are similar to deterministic TM except:
 - Zero or more alternatives can be defined for the transitions between states
 - the machine can choose the transition to be used.
 - The best of the possible transitions is taken into account to estimate complexity (the machine will guess it).
 - The transition function δ for the nondeterministic case is as follows:

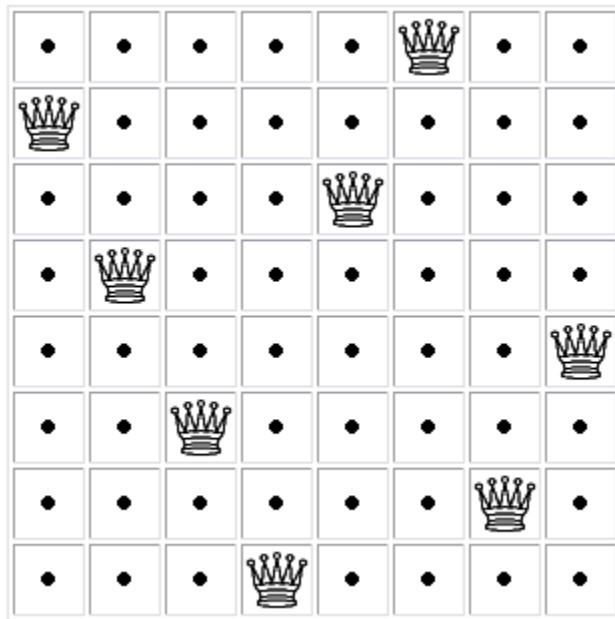
$$\delta : Q \times \Gamma \longrightarrow \mathcal{P}(Q \times \Gamma \times \{L, R\}).$$

is the set defined as parts of

$$\mathcal{P}(Q \times \Gamma \times \{L, R\}). \quad Q \times \Gamma \times \{L, R\}$$

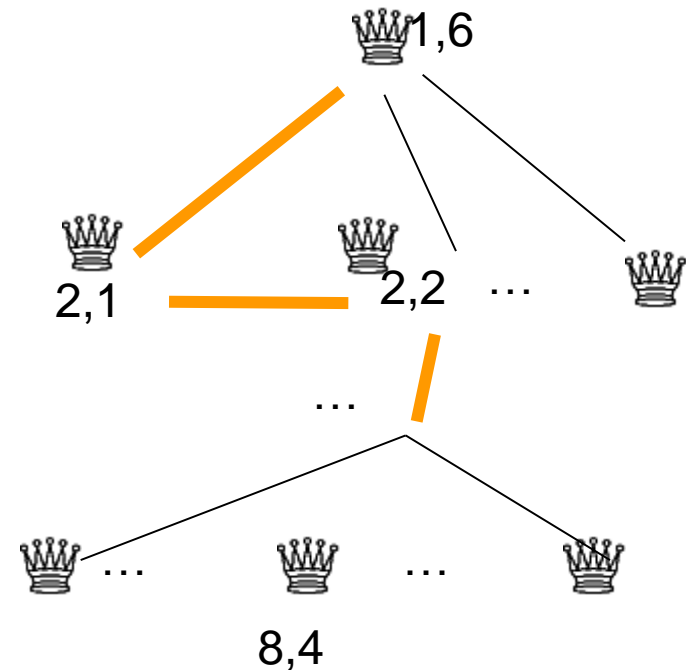


■ Eight queens puzzle



Place eight chess queens on an 8x8 chessboard so that no two queens attack each other.

Linear if the best solution is chosen everytime



http://en.wikipedia.org/wiki/Eight_queens_puzzle



Non-Deterministic Turing Machine

- Calculations in a ND-TM can be represented as a tree whose branches denote the possible possibilities for each transition.

If a “branch” reaches an acceptance state, the ND-TM accepts the input.

- Theorem:

Every Non-Determinist TM has an equivalent Determinist TM

- Therefore, Non-Determinist TM and Determinist TM are equivalent considering the computing capacity.





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■ Criterion (Theory of Computation):

- Decision Problem (Definition, instance, YES o NOT answer). Determine whether a “word” is included in the language or not.

■ Types:

- Decidable.
- Partially Decidable (Recognizable).
- Not Decidable.





■ Criterion (Computational Complexity Theory):

■ In terms of a Deterministic or Non-Deterministic TM solving the problem in a requiring polynomial time.

■ Types:

- Class P, NP, NP-Complete (CC Temporal).
- Class PSPACE, PSPACE-Complete, L, NL, NL-Complete (CC-Spatial).





Classification of Problems

- Let A be the universal language, TM a Turing Machine and $x \in A$:
Determine if x included in the language is a decision problem.

- Considering the *Theory of Computation*, a **Decision Problem can be:**
 1. **Decidable** (o with solution using an algorithm):
 - It is possible to find a device (TM) to solve the problem (i.e., for every x the TM returns a 1 or a 0 if x is included in the language or not).
 - In addition, the TM stops **for every element in A** (every input).





Classification of Problems

- Let A be the universal language, TM a Turing Machine and $x \in A$:
Determine if x included in the language is a decision problem.

- Considering the *Theory of Computation*, a **Decision Problem can be:**
 2. **Partially Decidable** (Recognizable):
 - It is possible to find a device (TM) to solve the problem (i.e., for every x the TM returns a 1 or a 0 if x is included in the language or not.
 - In addition, the TM stops **for every element in A included in the language recognized by the TM.**





Classification of Problems

- Let A be the universal language, TM a Turing Machine and $x \in A$:
Determine if x included in the language is a decision problem.

- Considering the *Theory of Computation*, a **Decision Problem can be:**
 3. **No Decidable:**
 - The problem is not decidable.





□ Criterion (Computational Complexity Theory):

- Categorize computational problems according to their behavior related to the Computational Complexity Theory:
 - Class **P** and Class **NP**,
 - whose solutions in a **polynomial time** are found considering **the determinism or non-determinism behavior respectively** of the TM.



□ Criterion (Computational Complexity Theory):

■ Class P:

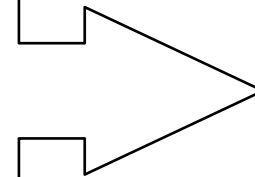
- Problems that a Deterministic TM can solve in a polynomial time related to the length of the input string.
 - Example: The language $L = \{a^n b^n c^n \mid n \geq 0\} \in P$ can be solved in a polynomial time using a Deterministic TM.
- Most of the usual problems are included in this class:
 - Order, Search, Operations with arrays...

□ Criterion (Computational Complexity Theory):

■ Class NP:

- Problems that a Non-Deterministic TM can solve in a polynomial time related to the length of the input string.

Given that every Deterministic TM is a particular case of a Non-Deterministic TM:



$$P \subseteq NP$$



□ Criterion (Computational Complexity Theory):

■ Class NP-Complete:

- A specific problem P is NP-Complete if:
 - It is P
 - All the other problems of NP can be reduced to it in polynomial time in a Deterministic-TM.

Reduce a problem.:

- * Transform a problem into another and then use the solution of the latter to obtain the solution of the former problem.





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Summary

- We have described the most common concepts and problems.
- There are problems which are more complex than other according to our computational model (Turing Machine)
- Although we cannot solve the most complex problems, we can know whether a specific problem is included in this category or not.
- The most complex problem is the NP-Complete.

