

	<p><b>UNIVERSIDAD CARLOS III DE MADRID</b></p> <p><b>FORMAL LANGUAGES AND AUTOMATA THEORY</b></p> <p><b>COMPUTER SCIENCE DEGREE. CONTINUOUS ASSESSMENT - PARTIAL 2</b></p> <p>Last name(s): _____</p> <p>First name: _____</p> <p>NIA: _____</p>
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**Duration: 45 minutes**

1. Indicate whether the following statements are true or false.

Correct Answer:+0.3

Wrong Answer:-0.3

No answer =0

Minimum Mark: 0

**Maximum Mark: 3 points**

	True	False
In a grammar, the word $\lambda$ can be included in the set of terminal symbols.		
A sentential form can include nonterminal symbols.		
There is an algorithm to determine whether a grammar is ambiguous or not.		
$A ::= aB$ is a rule in CNF.		
For every Type-1 grammar it is possible to find an equivalent Push-Down Automaton.		
Type-2 languages are also called context-free languages.		
Every type-2 grammar is also a type-3 grammar.		
If the number of sentences of a Language is infinite, then the grammar is recursive.		
Every grammar is GNF is a Type-3 grammar.		
$S \rightarrow AbbB \rightarrow BaA$ is a derivation of order 3.		

2. Indicate whether the following statements are true or false.

Correct Answer:+0.3      Wrong Answer:-0.3      No answer =0  
 Minimum Mark: 0      **Maximum Mark: 3 points**

	True	False
Given a grammar ( $\Sigma=\{0,1\}$ ) with the following production rules $\{ S ::= A 1 \mid 0 B \mid 0, A ::= A 1 \mid 1, B ::= \lambda \mid 0 S \}$ , it is a Type-3 grammar.		
The grammar with axiom S and production rules $P=\{S::=B0; B::=1 \mid C1; C::=B2; D::=S1\}$ is a recursive grammar.		
Given a grammar ( $G = \{\{0,1\},\{S,A,B\},S, P\}$ ) with the following production rules $P = \{ S ::= 1 S \mid 1 A \mid 0 B \mid 0 \mid \lambda, A ::= 1 A \mid 1, B ::= \lambda \mid 0 B \}$ , it can be transformed into a Type-3 left-linear grammar.		
$A::=A$ is a unit production or redenomination rule		
The production rule $AaC ::= AbabC$ , does not retain the context.		
Given the grammar, $G=( \{a,b\} , \{A,B,C, S\}, S, P=\{S::= aA; A::= aA \mid bB; B::= bB \mid b\}, , )$ , aabb is a sentence of the language.		
$CcdEA::=CCEb$ is a valid rule for a Type-1 grammar.		
Every Type-3 grammar has an equivalent NON-Deterministic FA.		
Every Type-3 grammar has an equivalent Deterministic FA.		
Given a grammar ( $G = \{\{0,1\},\{S,A,B\},S, P\}$ ) with the following production rules $P = \{ S ::= 1 S \mid 1 A \mid 0 B \mid 0 \mid \lambda, A ::= 1 A \mid 1, B ::= \lambda \mid 0 B \}$ , if the grammar is cleaned and becomes well-formed, then it would be a type-3 right-linear grammar.		

3. (4 points) Given the grammar  $G = (\{0, 1, 2\} , \{S, A, B, C, D, E\}, S, P)$

$P = \{(S ::= AABC), (A ::= \lambda \mid 1A0), (B ::= 1B \mid 1), (C ::= 1C1 \mid 0C0 \mid \lambda), (D ::= 0)\}$

- a) Which is the type of grammar in the Chomsky Hierarchy?
- b) Obtain an equivalent grammar in GNF. Explain in detail the process that you have followed.