

# Mathematical background

## CRYPTOGRAPHY AND COMPUTER SECURITY

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# OUTLINE

- 1. Mathematical background
  - Basic concepts
    - Congruence
    - Modular reduction
    - $Z_n$  set
  - Inverse computation
    - Fermat's theorem
    - Euler Totient Function
    - Euler's theorem
    - Inverse computation by means of Extended Euclidean Algorithm
  - Congruence equations
    - Powers of an integer
    - Primitive roots
    - Discrete logarithms

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# MATHEMATICAL BACKGROUND

- Basic concepts

- Let  $\mathbf{Z}$  be the set of integers and  $a, b, c \in \mathbf{Z}$

- $(\mathbf{Z}, +)$  qualifies as a group if:

$$a + b \in \mathbf{Z}$$

closure

$$a + (b + c) = (a + b) + c$$

associativity

$$a + 0 = a$$

identity element

$$a + (-a) = 0$$

inverse element

# MATHEMATICAL BACKGROUND

- $(\mathbf{Z}, +)$  qualifies as an abelian group if:

$$a + b = b + a$$

commutativity (abelian)

- $(\mathbf{Z}, +, \cdot)$  qualifies as a ring if  $(\mathbf{Z}, +)$  is an abelian group and:

$$a \cdot b \in \mathbf{Z}$$

closure

$$a \cdot (b \cdot c) = (a \cdot b) \cdot c$$

associativity

$$a \cdot 1 = a$$

identity element

$$a \cdot (b + c) = (a \cdot b) + (a \cdot c)$$

• distributivity over +

$(\mathbf{Z}, \cdot)$   
monoid

# MATHEMATICAL BACKGROUND

- $(\mathbf{Z}, +, \cdot)$  qualifies as a commutative ring if:

$$a \cdot b = b \cdot a$$

conmutativity

- $(\mathbf{Z}, +, \cdot)$  is a field if it is a commutative ring and possesses a multiplicative inverse:

$$a \cdot a^{-1} = 1$$

Inverse element ( $\cdot$ )

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# CONGRUENCE RELATIONSHIP

- For a positive integer  $n$ , two integers  $a$  and  $b$  are said to be **congruent modulo  $n$** , written:

$$a \equiv b \pmod{n} \Leftrightarrow (a - b) \text{ is an integer multiple of } n.$$

- The number  $n$  is called the **modulus** of the congruence.
- An equivalent definition is that both numbers have the same remainder when divided by  $n$ .
- $(\text{mod } n)$  operator maps all integers into the set  $\{0, 1, \dots, n-1\}$



# CONGRUENCE CLASSES

- We denote  $[a]$  to the set  $\{\dots, a-2n, a-n, a, a+n, a+2n, \dots\}$ , that is the congruence class of a modulo  $n$  (that is, for all  $x, y \in [a]$  regarding  $n$ ,  $x \equiv y \pmod{n}$  y  $a \in \{0, 1, \dots, n-1\}$ )

The congruence class  $[3]$  modulo 10 =

$$\begin{aligned} [3]_{10} &= \{\dots, -27, -17, -7, 3, 13, 23, 33, \dots\} = \\ &= \{\dots, 3 - 3 \cdot 10, 3 - 2 \cdot 10, 3 - 1 \cdot 10, 3 - 0 \cdot 10, 3 + 1 \cdot 10, 3 + 2 \cdot 10, 3 + 3 \cdot 10, \dots\} \\ &= -27 \equiv -17 \equiv -7 \equiv 3 \equiv 13 \equiv 23 \equiv 33 \pmod{10} \end{aligned}$$

The congruence class  $[7]$  modulo 11 =

$$\begin{aligned} [7]_{11} &= \{\dots, -26, -15, -4, 7, 18, 29, 40, \dots\} = \\ &= \{\dots, 7 - 3 \cdot 11, 7 - 2 \cdot 11, 7 - 1 \cdot 11, 7 - 0 \cdot 11, 7 + 1 \cdot 11, 7 + 2 \cdot 11, 7 + 3 \cdot 11, \dots\} \\ &= -26 \equiv -15 \equiv -4 \equiv 7 \equiv 18 \equiv 29 \equiv 40 \pmod{11} \end{aligned}$$

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# MODULAR REDUCTION

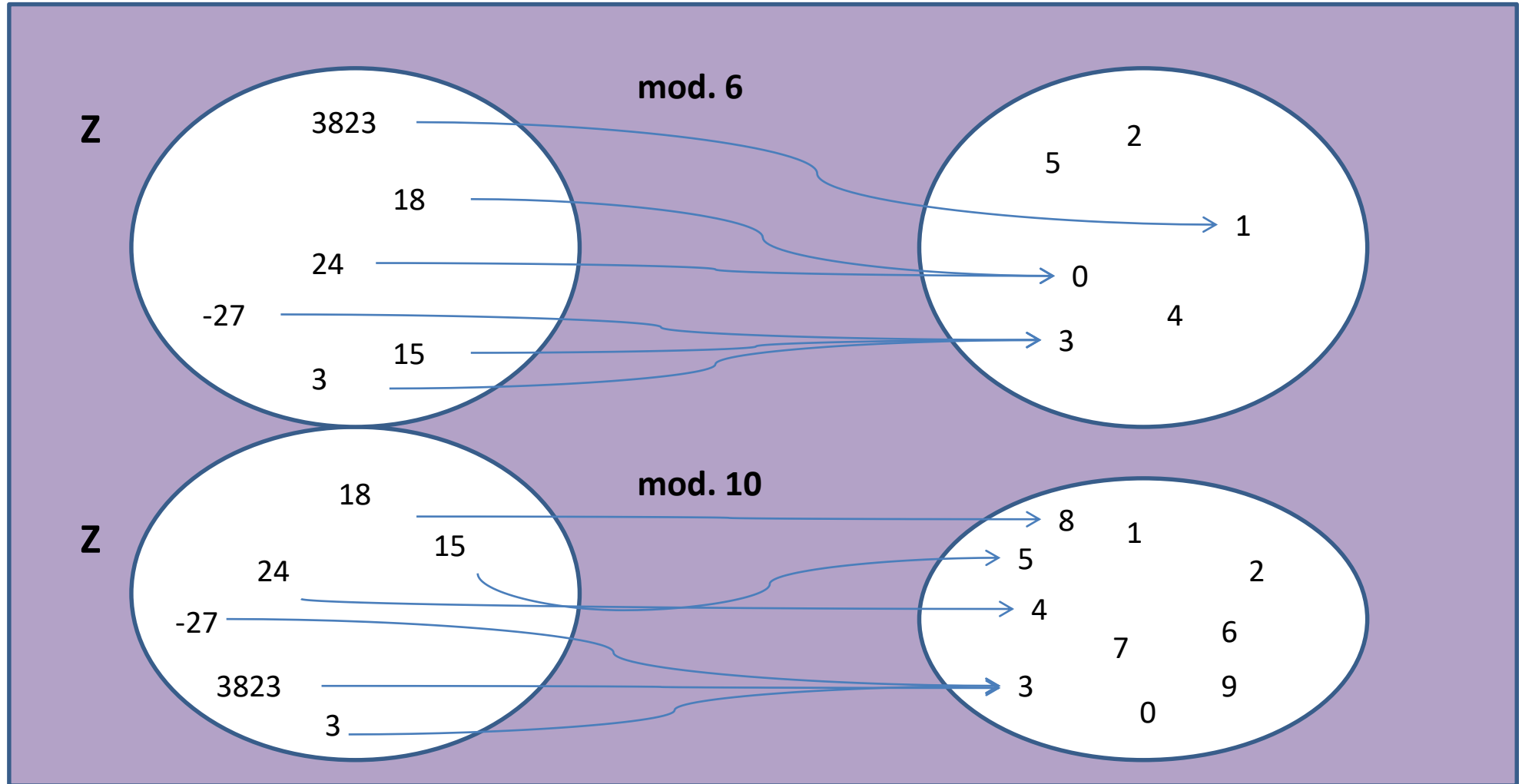
Given  $a, n \in \mathbf{Z}$  ( $n \neq 0$ ). Modular reduction (or modulo  $n$ ) is called to the function (represented by  $(\text{mod. } n)$ ) that applying to  $a$ , the result is  $r \in \mathbf{Z}^+ + \{0\} / r \in \{0, 1, \dots, n-1\}$  and  $a \equiv r \pmod{n}$

$$a \pmod{n} = r \Rightarrow a \equiv r \pmod{n} \text{ and } r \in \{0, 1, \dots, n-1\}$$

Note: “ $r$  is the remainder of the integer division of  $a$  between  $n$  (for  $a > 0$ )”

$26 \pmod{5} = 5 \cdot 5 + 1 \pmod{5} = 1$	$(1 < 5-1)$	$26 \equiv 1 \pmod{5}$
$30 \pmod{7} = 4 \cdot 7 + 2 \pmod{7} = 2$	$(2 < 7-1)$	$30 \equiv 2 \pmod{7}$
$11 \pmod{33} = 11$	$(11 < 33-1)$	
$256 \pmod{8} = 32 \cdot 8 + 0 \pmod{8} = 0$	$(0 < 8-1)$	$256 \equiv 0 \pmod{8}$
$-17 \pmod{12} \equiv -17 + 2 \cdot 12 = 7$	$(7 < 12-1)$	$-17 \equiv 7 \pmod{12}$

# MODULAR REDUCTION



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# $Z_n$ SET

$$- Z_n = \{ [a] / a \in Z \}$$

$Z_n$  is the set of congruence classes regarding a modulo  $n$

If  $n \neq 0$ ,  $Z_n = \{ [0], [1], \dots, [n-1] \}$  or simplified

$Z_n = \{ 0, 1, \dots, n-1 \}$  (“**the ring of integers modulo  $n$** ”)

$$Z_8 = \{0, 1, 2, 3, 4, 5, 6, 7\}$$

$$Z_{31} = \{0, 1, 2, 3, 4, 5, 6, 7, \dots, 26, 27, 28, 29, 30\}$$

# $Z_n$ SET

- **Operations ( $+_n, \cdot_n$ )**

- Addition and multiplication is defined for  $Z_n$

For  $[a], [b], [c] \in Z_n$        $+_n : [a] +_n [b] = [a+b]$        $\cdot_n : [a] \cdot_n [b] = [a \cdot b]$

$Z_8 :$

$$[6] +_8 [7] = (6 + 7) \text{ mod. } 8 = 13 \text{ mod. } 8 = [5]$$

$Z_{31} :$

$$[3] \cdot_{31} [11] = (3 \cdot 11) \text{ mod. } 31 = 33 \text{ mod. } 31 = [2]$$

# $Z_n$ SET

## – $Z_n$ - properties regarding $(+_n, \cdot_n)$

For  $n \neq 0$ ,  $Z_n$  is a computative ring regarding  $(+_n, \cdot_n)$

$$[a] +_n [b] \in Z_n$$

$$[a] \cdot_n [b] \in Z_n$$

$$[a] +_n ([b] +_n [c]) = ([a] +_n [b]) +_n [c]$$

$$[a] \cdot_n ([b] \cdot_n [c]) = ([a] \cdot_n [b]) \cdot_n [c]$$

$$[a] +_n 0 = [a]$$

$$[a] \cdot_n 1 = [a]$$

$$[a] +_n (-[a]) = 0$$

$$[a] \cdot_n ([b] +_n [c]) = ([a] \cdot_n [b]) +_n ([a] \cdot_n [c])$$

$$[a] +_n [b] = [b] +_n [a]$$

$$[a] \cdot_n [b] = [b] \cdot_n [a]$$



# $Z_n$ SET

- ▶ Homomorphism relationship with  $Z$  ring of integers
- ▶ Function “Modular reduction” is an homomorphism between  $Z$  (ring of integers) and  $Z_n$  (ring of integers modulo  $n$ )  $\leftrightarrow$  It is verified that:

Given  $a, b \in Z, f(a), f(b) \in Z_n$ , with  $f: Z \rightarrow Z_n$  ( $f =$  Modular reduction)

$$f(a + b) = f(a) +_n f(b) \quad ; \quad f(a \cdot b) = f(a) \cdot_n f(b)$$

- ▶ “Consequences” (Fundamental principles of modular arithmetic):

- $(a + b) \pmod{n} = a \pmod{n} +_n b \pmod{n} = (a \pmod{n} + b \pmod{n}) \pmod{n}$

$$[ (a + b) ] = [a] +_n [b] = [ [a] + [b] ]$$

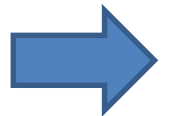
- $(a \cdot b) \pmod{n} = a \pmod{n} \cdot_n b \pmod{n} = (a \pmod{n} \cdot b \pmod{n}) \pmod{n}$

$$[ (a \cdot b) ] = [a] \cdot_n [b] = [ [a] \cdot [b] ]$$

- $(a \cdot (b+c)) \pmod{n} = ((a \pmod{n} \cdot_n b \pmod{n}) +_n (a \pmod{n} \cdot_n c \pmod{n})) \pmod{n}$   
 $= ((a \cdot b) \pmod{n} + (a \cdot c) \pmod{n}) \pmod{n}$

$$[ (a \cdot (b + c)) ] = [a] \cdot_n ( [b] +_n [c] ) = [ [a] \cdot ( [b] + [c] ) ]$$

examples



# MATH FUNDAMENTALS

$$(a + b)(\text{mod. } n) = a (\text{mod. } n) +_n b (\text{mod. } n) = (a(\text{mod. } n) + b(\text{mod. } n))(\text{mod. } n)$$

$$[ (a + b) ] = [a] +_n [b] = [ [a] + [b] ]$$

Example:

$$\begin{aligned} (3 + 8) (\text{mod. } 5) &= 3 (\text{mod. } 5) +_n 8 (\text{mod. } 5) = \\ &= (3 (\text{mod. } 5) + 8 (\text{mod. } 5)) \text{mod. } 5 = \\ &= (3 + 3) \text{mod. } 5 = 6 \text{mod. } 5 = \mathbf{1} \end{aligned}$$

$$(3 + 8) (\text{mod. } 5) = 11 (\text{mod. } 5) = 1$$

# MATH FUNDAMENTALS

$$(a \cdot b)(\text{mod. } n) = a(\text{mod. } n) \cdot_n b(\text{mod. } n) = (a(\text{mod. } n) \cdot b(\text{mod. } n))(\text{mod. } n)$$

$$[ (a \cdot b) ] = [a] \cdot_n [b] = [ [a] \cdot [b] ]$$

Example:

$$\begin{aligned} (3 \cdot 8)(\text{mod. } 5) &= 3(\text{mod. } 5) \cdot_n 8(\text{mod. } 5) = (3(\text{mod. } 5) \cdot 8(\text{mod. } 5)) \text{mod. } 5 = \\ &= (3 \cdot 3) \text{mod. } 5 = 9 \text{mod. } 5 = \mathbf{4} \end{aligned}$$

$$(3 \cdot 8)(\text{mod. } 5) = 24(\text{mod. } 5) = 4$$

$$\begin{aligned} 7^4(\text{mod. } 5) &= (7 \cdot 7 \cdot 7 \cdot 7)(\text{mod. } 5) = \\ &= (7(\text{mod. } 5) \cdot 7(\text{mod. } 5) \cdot 7(\text{mod. } 5) \cdot 7(\text{mod. } 5)) \text{mod. } 5 \Rightarrow \\ &= (2 \cdot 2 \cdot 2 \cdot 2) \text{mod. } 5 = 2^4(\text{mod. } 5) = 16 \text{mod. } 5 = \mathbf{1} \end{aligned}$$

$$7^4(\text{mod. } 5) = 2401(\text{mod. } 5) = 1$$

# MATH FUNDAMENTALS

$$\begin{aligned} (a \cdot (b+c))(\text{mod. } n) &= ((a(\text{mod. } n) \cdot_n b(\text{mod. } n)) +_n (a(\text{mod. } n) \cdot_n c(\text{mod. } n))) = \\ &= ((a \cdot b)(\text{mod. } n) + (a \cdot c)(\text{mod. } n))(\text{mod. } n) \\ [ (a \cdot (b + c)) ] &= [a] \cdot_n ( [b] +_n [c] ) = [ [a] \cdot ( [b] + [c] ) ] \end{aligned}$$

Example:

$$\begin{aligned} (3 \cdot (8+4))(\text{mod. } 5) &= ((3(\text{mod. } 5) \cdot_n 8(\text{mod. } 5)) +_n (3(\text{mod. } 5) \cdot_n 4(\text{mod. } 5))) = \\ &= ((3 \cdot 8)(\text{mod. } 5) + (3 \cdot 4)(\text{mod. } 5))(\text{mod. } 5) = \\ &= ((3 \cdot 3)(\text{mod. } 5) + (3 \cdot 4)(\text{mod. } 5))(\text{mod. } 5) = \\ &= (9 (\text{mod. } 5) + 12 (\text{mod. } 5)) (\text{mod. } 5) = (4 + 2) (\text{mod. } 5) = \\ &= 6 \text{ mod. } 5 = \mathbf{1} \end{aligned}$$

$$(3 \cdot (8+4))(\text{mod. } 5) = 36 (\text{mod. } 5) = 1$$

# MATH FUNDAMENTALS

More examples:

$$(23 + 4)(\text{mod. } 5) = 2$$

$$2^9 (\text{mod. } 5) = 2$$

$$(3 + 8) \cdot 5 (\text{mod. } 5) = 0$$

$$(41 + 1001) \cdot 999 (\text{mod. } 5) = 3$$

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# FERMAT'S THEOREM

If  $p$  is a prime,  $\forall a \in \mathbf{Z} / \text{g.c.d.}(a, p) = 1$  then:

$$a^{p-1} \pmod{p} = 1$$

- $a * a^{p-2} = 1 \pmod{p}$ , so  $a^{p-2}$  is the inverse of  $a \pmod{p}$
- Also known as Fermat's Little Theorem
- Also:  $a^p = a \pmod{p}$
- Useful in public key and primality testing

# COMPUTING THE INVERSE. FERMAT'S THEOREM

- Compute:  $2x \bmod 7 = 1$

## **Solution:**

$a=2$ ,  $p=7$  prime,  $\text{g.c.d.}(2,7)=1$ , applying Fermat:

$$x = 2^{p-2} \bmod 7 \Leftrightarrow x = 2^{7-2} \bmod 7 \Leftrightarrow x = 2^5 \bmod 7 \Leftrightarrow$$

$$x = 2^3 \cdot 2^2 \bmod 7 \Leftrightarrow x = 4 \bmod 7$$

- Compute:  $35x \bmod 3 = 1$

## **Solution**

$a=35$ ,  $p=3$  prime,  $\text{g.c.d.}(35,3)=1$ ,

$$35x \bmod 3 = (35 \bmod 3) (x \bmod 3) \bmod 3 = 2x \bmod 3 = 1$$

applying Fermat:  $x = 2^{p-2} \bmod 3 \Leftrightarrow x = 2^{3-2} \bmod 3 \Leftrightarrow x = 2$



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# EULER TOTIENT FUNCTION $\Phi(N)$

- **Reduced set of residues** ( $Z_n^*$ ) is composed of numbers (residues) which are relatively prime to  $n$ 
  - e.g. for  $n=10$ ,
  - reduced set of residues is  $Z_n^* = \{1,3,7,9\}$
- Number of elements (order) in reduced set of residues is called the **Euler Totient Function  $\Phi(n)$**

# EULER TOTIENT FUNCTION $\Phi(N)$

- **Computing  $\Phi(n)$**

- If  $p$  is prime

$$\Phi(p) = p-1$$

- If  $p$  is prime and  $k \in \mathbf{Z}^+$ :

$$\Phi(p^k) = p^k - p^{k-1} = p^{k-1}(p-1)$$

- If  $p$  and  $q$  are relatively prime:

$$\Phi(pq) = \Phi(p) \Phi(q)$$

- If  $n = \prod p_i^{k_i} / \forall i p_i$  is prime,  $k_i \in \mathbf{Z}^+$ :

$$\Phi(n) = \prod p_i^{k_i-1} (p_i - 1)$$

# EULER TOTIENT FUNCTION $\Phi(N)$

- Examples

$$\Phi(37) = 36$$

$$\Phi(21) = (3-1) \cdot (7-1) = 2 \cdot 6 = 12$$

$$\Phi(172) = ?$$

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# EULER'S THEOREM

- A generalization of Fermat's theorem

$\forall a, n \in \mathbf{Z} (n \neq 0) / \text{g.c.d.}(a, n) = 1:$

$$a^{\Phi(n)} \pmod{n} = 1$$

so:

$$a a^{\Phi(n)-1} \pmod{n} = 1$$

$$a^{-1} = a^{\Phi(n)-1} \pmod{n}$$

- If  $n$  is prime (denoted as  $p$ ):

$$a^{-1} = a^{p-2} \pmod{p}$$

# EULER'S THEOREM

- Examples

- $3x \bmod 10 = 1$

$$a=3; n=10; \Phi(10)=4;$$

$$\begin{aligned} x &= 3^{4-1} \bmod 10 = 3^3 \bmod 10 = 3 \cdot 3^2 \bmod 10 \\ &= 3 \cdot (-1) \bmod 10 = -3 \bmod 10 = 7 \end{aligned}$$

- $2x \bmod 11 = 1$

$$a=2; n=11; \Phi(11)=10;$$

$$\begin{aligned} x &= 2^{10-1} \bmod 11 = (2^3)^3 \bmod 11 = (-3)^3 \bmod 11 = \\ &= (-2) \cdot (-3) \bmod 11 = 6 \end{aligned}$$

# COMPUTING THE INVERSE. EULER'S THEOREM. EXAMPLES:

- Compute:  $37x \bmod 10 = 1$

## Solution

$$a=37, n=10, \text{g.c.d.}(37,10)=1, 7x \bmod 10 = 1$$

$$\text{applying Euler: } x = 7^{\Phi(10)-1} \bmod 10$$

$$10 = 2 * 5, \Phi(10) = \Phi(2) * \Phi(5) = 1 * 4 = 4$$

$$x = 7^{4-1} \bmod 10 \Leftrightarrow x = 7^3 \bmod 10 \Leftrightarrow x = (-1) 7 \bmod 10 \Leftrightarrow$$

$$x = -7 \bmod 10 \Leftrightarrow x = 3 \bmod 10 = 3$$



# COMPUTING THE INVERSE. EULER'S THEOREM. EXAMPLES:

- Compute:  $17x \bmod 12 = 1$

## Solution

$$a=17, n=12, \text{g.c.d.}(17,12)=1, 5x \bmod 12 = 1$$

$$\text{applying Euler: } x = 5^{\Phi(12)-1} \bmod 12$$

$$12 = 2^2 * 3, \Phi(12) = \Phi(2^2) * \Phi(3) = 2^{2-1} * (2-1) * 2 = 4$$

$$x = 5^{4-1} \bmod 12 \Leftrightarrow x = 5^3 \bmod 12 \Leftrightarrow x = 13 \cdot 5 \bmod 12$$

$$\Leftrightarrow x = 5$$

# COMPUTING THE INVERSE.

## EXAMPLES:

- Compute:  $37x \bmod 41 = 1$

### **Solution**

$a=37$ ,  $n=41$ ,  $\text{g.c.d.}(37,41)=1$ ,  
applying Euler:  $x = 37^{\Phi(41) - 1} \bmod 41$

...

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# EUCLIDEAN ALGORITHM

- Determines the GCD of two integers.
- Its major significance is that it does not require factoring the two integers
- This technique can be used to compute the inverse of a number in modular arithmetic

# EUCLIDEAN ALGORITHM

- Example: Compute the  $\text{gcd}(1547, 560)$

	2	1	3	4	1	3
1547	560	427	133	28	21	7
427	133	28	21	<b>7</b>	<b>0</b>	

# EUCLIDEAN ALGORITHM

- $1547 = 2 * 560 + 427$
- $560 = 1 * 427 + 133$
- $427 = 3 * 133 + 28$
- $133 = 4 * 28 + 21$
- $28 = 1 * 21 + 7$
- $21 = 3 * 7 + 0$
- $\text{gcd}(1547, 560) = 7$

# EXTENDED EUCLIDEAN ALGORITHM

- Computing inverses

If  $\gcd(a,n)=1$

	$c_1$	$c_2$	...	...	...	$c_n$	$r_{n-1}$
$n$	$a$	$r_1$	$r_2$	...	...	$r_{n-1}$	<b>1</b>
$r_1$	$r_2$	$r_3$	...	...	<b>1</b>	<b>0</b>	

$$\begin{aligned}
 n &= c_1 a + r_1 \\
 a &= c_2 r_1 + r_2 \\
 r_1 &= c_3 r_2 + r_3 \\
 &\dots \\
 &\dots \\
 r_{n-2} &= c_n r_{n-1} + 1 \\
 r_{n-1} &= c_{n+1} + 0
 \end{aligned}$$

Substituting:

$$1 = k_1 a + k_2 n$$

Reducing modulo  $n$ :

$$1 = k_1 a \pmod{n}$$

Then

$$k_1 = a^{-1} \pmod{n}$$

# EXTENDED EUCLIDEAN ALGORITHM

- $37x \bmod 41 = 1$

	1	9	4
41	37	4	1
4	1	0	

$$\begin{aligned}n &= c_1 a + r_1 \Rightarrow \\41 &= 1 \cdot 37 + 4 \\37 &= 9 \cdot 4 + 1\end{aligned}$$

$$1 = 37 - 9 \cdot 4 = 37 - 9(41 - 37)$$

$$1 = 37 - 9 \cdot 41 + 9 \cdot 37 = 37 \cdot 10 - 9 \cdot 41$$

$$1 = 37 \cdot 10 \bmod 41$$

$$\mathbf{x = 10 \bmod 41}$$



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# LINEAR CONGRUENCE EQUATIONS

$$a x \equiv b \pmod{n}$$



$$a x + n k = b$$

- If  $\text{g.c.d.}(a,n)=1$ , the eq. has a unique solution

$$a y \equiv 1 \pmod{n} \{y = a^{-1} \pmod{n}\}; x \equiv b y \pmod{n}$$

- If  $\text{g.c.d.}(a,n) = d \neq 1, d \mid b, \exists d$  solutions between 0 and  $d-1$

$$x \equiv (b/d) y + j (n/d) \pmod{n} \quad j \in \{0, d-1\}$$

$$(a/d) y \pmod{n/d} = 1$$

- Else, there is no solution

# POWERS OF AN INTEGER, modulo $n$

- $a, n \in \mathbb{Z}$

$$a^0 \bmod n$$

$$a^1 \bmod n$$

$$a^2 \bmod n$$

...

$$a^g \bmod n$$

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# PRIMITIVE ROOTS. GENERATOR

- From Euler's theorem  $a^{\varphi(n)} \bmod n = 1$
- Consider  $a^g = 1 \pmod n$ ,  $\text{GCD}(a, n) = 1$ 
  - At least  $g = \varphi(n)$ , but may be smaller
  - **Order** of  $a \bmod n$  is the least  $g$  that holds eq.
  - if smallest  $g$  is  $g = \varphi(n)$  then  $a$  is called a **primitive root**

# PRIMITIVE ROOTS. GENERATOR

- Let  $p$  be a prime. If  $a \in \mathbf{Z}^+$  and  $\text{g.c.d.}(a, p)=1$ , then

the order of  $a$ ,  $\text{ord}(a)$ , modulo  $p$  is a divisor of  $p - 1$

- If  $p$  is prime and  $\text{ord}(a) = p - 1$ , then successive powers of  $a$  "generate" the group  $\text{mod } p$  ( $\mathbf{Z}_p^*$ )

There are  $\Phi(p-1)$  primitive roots modulo  $p$

# PRIMITIVE ROOTS. GENERATOR

- Compute the primitive roots of 7 ( $\mathbb{Z}_7^*$ )

$$2^0 \equiv 1, 2^1 \equiv 2, 2^2 \equiv 4, 2^3 \equiv 1 \qquad g = 3$$

$$3^0 \equiv 1, 3^1 \equiv 3, 3^2 \equiv 2, 3^3 \equiv 6, 3^4 \equiv 4, 3^5 \equiv 5, 3^6 \equiv 1, \qquad g = 6$$

$$4^0 \equiv 1, 4^1 \equiv 4, 4^2 \equiv 2, 4^3 \equiv 1 \qquad g = 3$$

$$5^0 \equiv 1, 5^1 \equiv 5, 5^2 \equiv 4, 5^3 \equiv 6, 5^4 \equiv 2, 5^5 \equiv 3, 5^6 \equiv 1 \qquad g = 6$$

$$6^0 \equiv 1, 6^1 \equiv 6, 6^2 \equiv 1 \qquad g = 2$$

- Solution: 3 and 5

# OUTLINE

- 1. Mathematical background
  - Basic concepts
    - Congruence
    - Modular reduction
    - $Z_n$  set
  - Inverse computation
    - Fermat's theorem
    - Euler Totient Function
    - Euler's theorem
    - Inverse computation by means of Extended Euclidean Algorithm
  - Congruence equations
    - Powers of an integer
    - Primitive roots
    - Discrete logarithms



# DISCRETE LOGARITHM

- the inverse problem to exponentiation is to find the **discrete logarithm** of a number modulo  $n$
- that is to find  $x$  such that  $b = a^x \pmod{n}$
- this is written as  $x = \log_a b \pmod{n}$
  
- If  $a$  is a primitive root then it always exists, otherwise it may not, eg.
  - $x = \log_3 4 \pmod{13}$  has no answer
  - $x = \log_2 3 \pmod{13} = 4$  by trying successive powers

# DISCRETE LOGARITHM

a	a <sup>2</sup>	a <sup>3</sup>	a <sup>4</sup>	a <sup>5</sup>	a <sup>6</sup>	a <sup>7</sup>	a <sup>8</sup>	a <sup>9</sup>	a <sup>10</sup>	a <sup>11</sup>	a <sup>12</sup>	a <sup>13</sup>	a <sup>14</sup>	a <sup>15</sup>	a <sup>16</sup>	a <sup>17</sup>	a <sup>18</sup>
3	9	8	5	15	7	2	6	18	16	10	11	14	4	12	17	13	1
4	16	7	9	17	11	6	5	1	4	16	7	9	17	11	6	5	1
7	11	1	7	11	1	7	11	1	7	11	1	7	11	1	7	11	1

- $x = \log_3 7 \pmod{19} = 6$  (3 is primitive root modulo 19)
- $x = \log_4 3 \pmod{19}$  –There is not solution, 4 is not primitive root ( $4^{ix} = 3 \pmod{19}$ ) but
- $x = \log_4 9 \pmod{19}$  –There is solution but it is not unique  $x = \{4, 13\}$

# CRYPTOGRAPHY AND COMPUTER SECURITY

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