



“Symmetric Encryption: Block ciphers”

Proposed exercises

Exercise 1 :

Assume the following DES key:

10000101 10100100 10001111 10001111 10000101 10100100 10001111 10001111.

- Compute the first internal subkey generated by the algorithm to encrypt a cleartext.
- Compute L1 y R1 for the following cleartext: **10101010 10101010 10101010 10101010 10101010 10101010 10101010 10101010**

Exercise 2:

Consider a DES cipher in CBC mode, and the following data:

The cleartext message M = **10101010 10101010 10101010 10101010 10101010 10101010 10101010 10101010 01010101 01010101 01010101 01010101 01010101 01010101 01010101 01010101**

The initial value for the registry C_0 = **11111111 00000000 11111111 00000000 11111111 00000000 11111111 00000000**

- Compute the input value to the S-BOX in the first iteration, assuming that there the IP permutation is not performed, and the first internal subkey is k_1 = **000000 111111 000000 111111 000000 111111 000000 111111**.
- Assuming that, after the first iteration of the encryption process, the output of the cipher is C_1 = **01010101 01010101 01010101 01010101 01010101 01010101 01010101 01010101**, compute the input to the block cipher in the next iteration.
- Suppose that C_1 is sent over a communication line, and that there is a transmission error which affects 2 bits of this block. Explain and reason how this error would affect the decryption of the message.

Exercise 3:

We know that a user's DES key is composed by 8 symbols from an alphabet of 26 letters.

Considering that the time needed to test one single key is 1 microsecond, calculate:

- The time needed to break a cryptogram.
- The time needed, assuming an alphabet that also includes digits.

Exercise 4:

Given the following intermediate AES state 3 (i.e., the output of the ShiftRows function),

calculate the byte from row 1, column 0 (consider that the byte D4 is in position r0,c0):

D4	E0	B8	1E
BF	B4	41	27
5D	52	11	98
30	AE	F1	E5

Exercise 5:

AES SubByte function is a non-linear substitution which is applied independently to every byte within the status matrix (intermediate status 1). For this purpose, the S- BOX substitution table is employed. This table is build using two different transformations

- First: Calculate the multiplicative inverse of that byte with respect to the polynomial

$$m(x) = x^8 + x^4 + x^3 + x + 1$$

- Second: Apply the following transformation:

$$\begin{pmatrix} b'_0 \\ b'_1 \\ b'_2 \\ b'_3 \\ b'_4 \\ b'_5 \\ b'_6 \\ b'_7 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 0 & 0 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 & 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 \end{pmatrix} \begin{pmatrix} b_0 \\ b_1 \\ b_2 \\ b_3 \\ b_4 \\ b_5 \\ b_6 \\ b_7 \end{pmatrix} + \begin{pmatrix} 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 0 \end{pmatrix}$$

where x_i bits are parts of the result of the first transformation and y_i are the resulting bits of the second transformation (note: subindex 0 indicates

the least significant bit)

Suppose the byte A=10001000. Get the resulting byte using the transformations previously described. Check the resulting value using the S-BOX table below.

		y															
		0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
x	0	63	7c	77	7b	f2	6b	6f	c5	30	01	67	2b	fe	d7	ab	76
	1	ca	82	c9	7d	fa	59	47	f0	ad	d4	a2	af	9c	a4	72	c0
	2	b7	fd	93	26	36	3f	f7	cc	34	a5	e5	f1	71	d8	31	15
	3	04	c7	23	c3	18	96	05	9a	07	12	80	e2	eb	27	b2	75
	4	09	83	2c	1a	1b	6e	5a	a0	52	3b	d6	b3	29	e3	2f	84
	5	53	d1	00	ed	20	fc	b1	5b	6a	cb	be	39	4a	4c	58	cf
	6	d0	ef	aa	fb	43	4d	33	85	45	f9	02	7f	50	3c	9f	a8
	7	51	a3	40	8f	92	9d	38	f5	bc	b6	da	21	10	ff	f3	d2
	8	cd	0c	13	ec	5f	97	44	17	c4	a7	7e	3d	64	5d	19	73
	9	60	81	4f	dc	22	2a	90	88	46	ee	b8	14	de	5e	0b	db
	a	e0	32	3a	0a	49	06	24	5c	c2	d3	ac	62	91	95	e4	79
	b	e7	c8	37	6d	8d	d5	4e	a9	6c	56	f4	ea	65	7a	ae	08
	c	ba	78	25	2e	1c	a6	b4	c6	e8	dd	74	1f	4b	bd	8b	8a
	d	70	3e	b5	66	48	03	f6	0e	61	35	57	b9	86	c1	1d	9e
	e	e1	f8	98	11	69	d9	8e	94	9b	1e	87	e9	ce	55	28	df
	f	8c	a1	89	0d	bf	e6	42	68	41	99	2d	0f	b0	54	bb	16

Exercise 6:

The following matrix is the input matrix to the ByteSub function::

$$\begin{pmatrix} 09 & 93 & 19 & 27 \\ AE & 52 & 11 & 9D \\ 19 & 21 & A5 & 9C \\ A9 & CC & 33 & 30 \end{pmatrix}$$

Recall that the ByteSub transformation is based on the following table:

	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
0	63	7c	77	7b	f2	6b	6f	c5	30	01	67	2b	fe	d7	ab	76
1	ca	82	c9	7d	fa	59	47	f0	ad	d4	a2	af	9c	a4	72	c0
2	b7	fd	93	26	36	3f	f7	cc	34	a5	e5	f1	71	d8	31	15
3	04	c7	23	c3	18	96	05	9a	07	12	80	e2	eb	27	b2	75
4	09	83	2c	1a	1b	6e	5a	a0	52	3b	d6	b3	29	e3	2f	84
5	53	d1	00	ed	20	fc	b1	5b	6a	cb	be	39	4a	4c	58	cf
6	d0	ef	aa	fb	43	4d	33	85	45	f9	02	7f	50	3c	9f	a8
7	51	a3	40	8f	92	9d	38	f5	bc	b6	da	21	10	ff	f3	d2
8	cd	0c	13	ec	5f	97	44	17	c4	a7	7e	3d	64	5d	19	73
9	60	81	4f	dc	22	2a	90	88	46	ee	b8	14	de	5e	0b	db
a	e0	32	3a	0a	49	06	24	5c	c2	d3	ac	62	91	95	e4	79
b	e7	c8	37	6d	8d	d5	4e	a9	6c	56	f4	ea	65	7a	ae	08
c	ba	78	25	2e	1c	a6	b4	c6	e8	dd	74	1f	4b	bd	8b	8a
d	70	3e	b5	66	48	03	f6	0e	61	35	57	b9	86	c1	1d	9e
e	e1	f8	98	11	69	d9	8e	94	9b	1e	87	e9	ce	55	28	df
f	8c	a1	89	0d	bf	e6	42	68	41	99	2d	0f	b0	54	bb	16

- Calculate the output status matrix of the ByteSub function.
- After this function, the ShiftRow function is applied in AES. Calculate the output status matrix of the ShiftRow function
- Afterwards, the MixColumns function is applied. It is based on this transformation:

$$\begin{pmatrix} S'_{0,C} \\ S'_{1,C} \\ S'_{2,C} \\ S'_{3,C} \end{pmatrix} = \begin{pmatrix} 02 & 03 & 01 & 01 \\ 01 & 02 & 03 & 01 \\ 01 & 01 & 02 & 03 \\ 03 & 01 & 01 & 02 \end{pmatrix} \begin{pmatrix} S_{0,C} \\ S_{1,C} \\ S_{2,C} \\ S_{3,C} \end{pmatrix}$$

Taking as the input status matrix the one calculated previously, calculate the transformation of the column number 0 of that matrix